CRASH/SAFE: Clean-slate Co-design of a Secure Host Architecture

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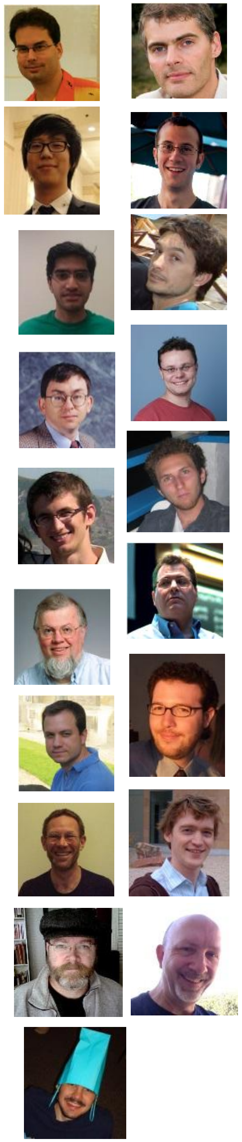
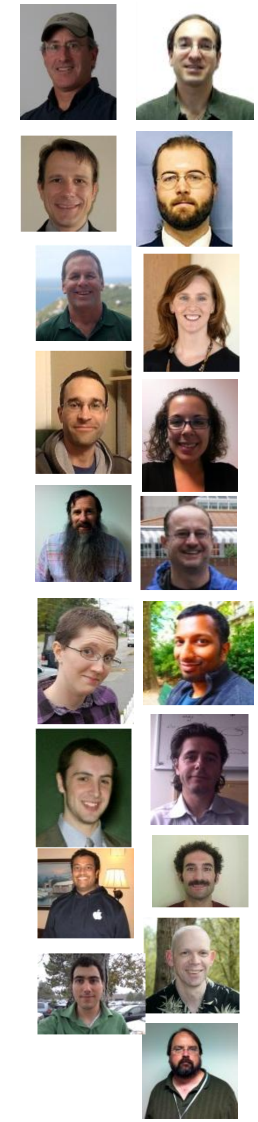
CRASH/SAFE project



- Academic partners (16):
  - **University of Pennsylvania** (11)
  - **Harvard University** (4)
  - **Northeastern University** (1)
- Industrial partners (24):
  - **BAE systems** (21) + **Clozure** (3)

40!

- Funded by DARPA
  - **C**lean-Slate Design of **R**esilient, **A**daptive, **S**ecure **H**osts



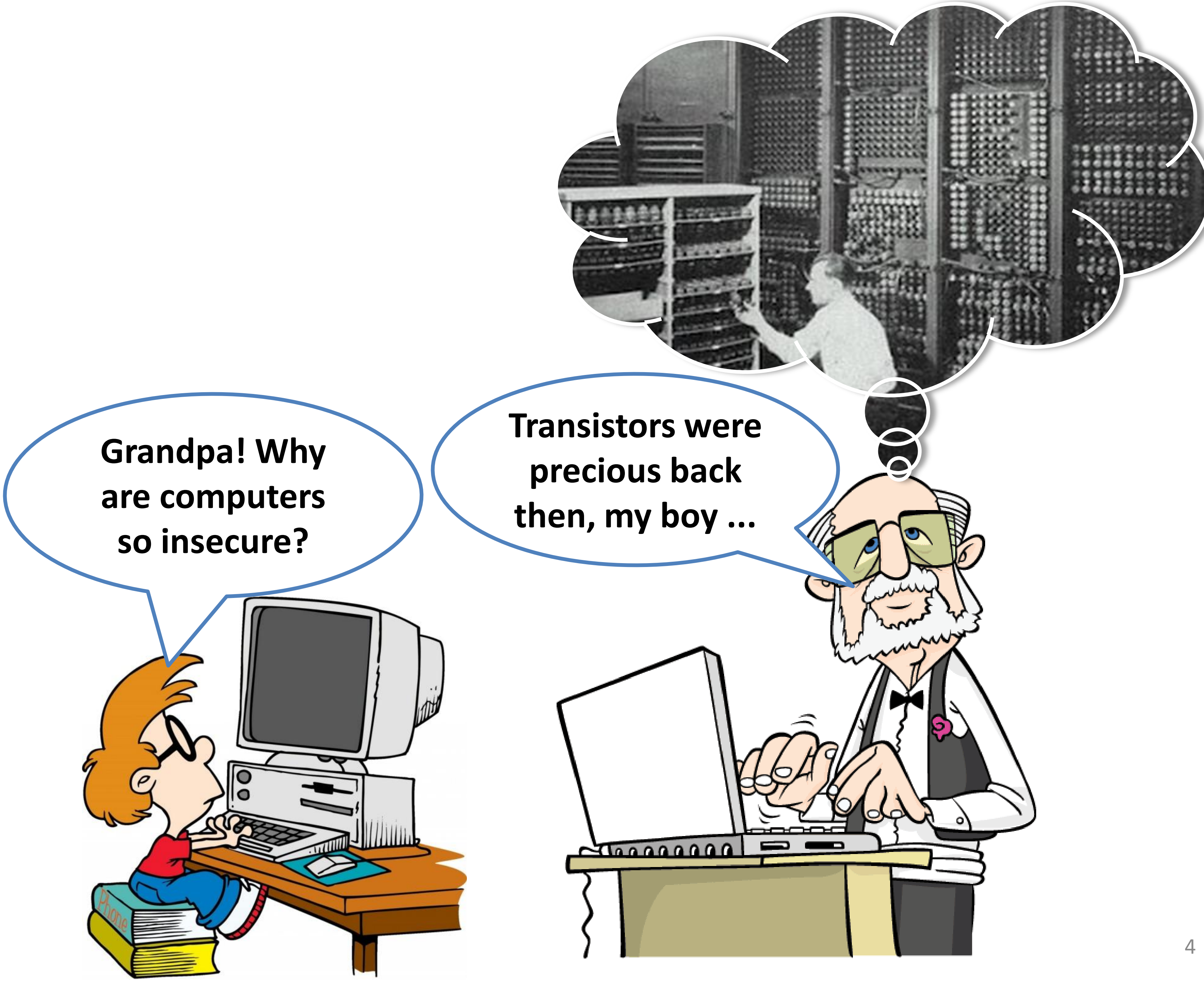
Clean-slate co-design of net host

Primary goal:
design and implement a significantly more secure architecture, without backwards compatibility concerns

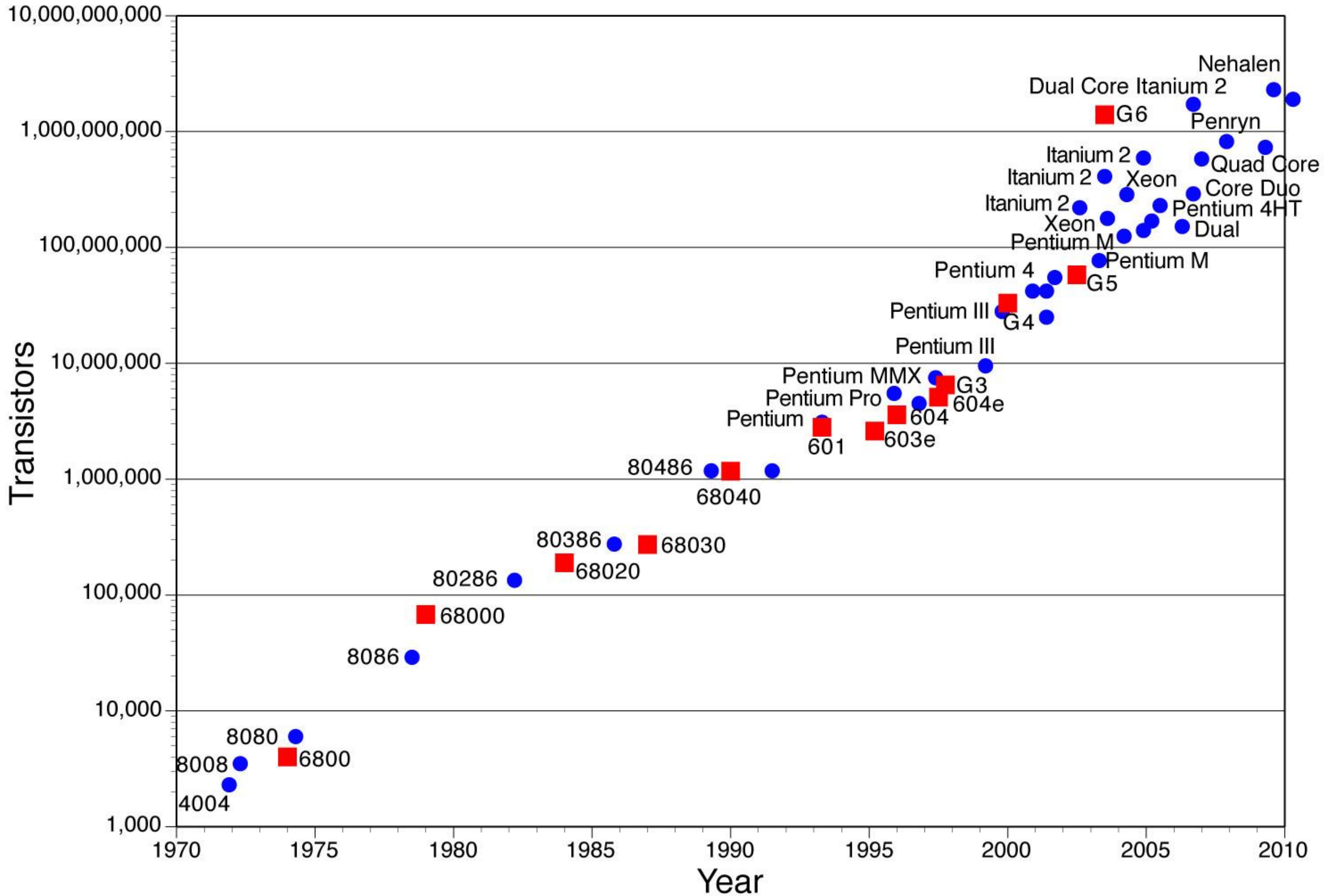
Secondary goal:
verify that it's secure (whatever that means)

New stack:

- language
- runtime
- hardware


Grandpa! Why are computers so insecure?
Transistors were precious back then, my boy ...

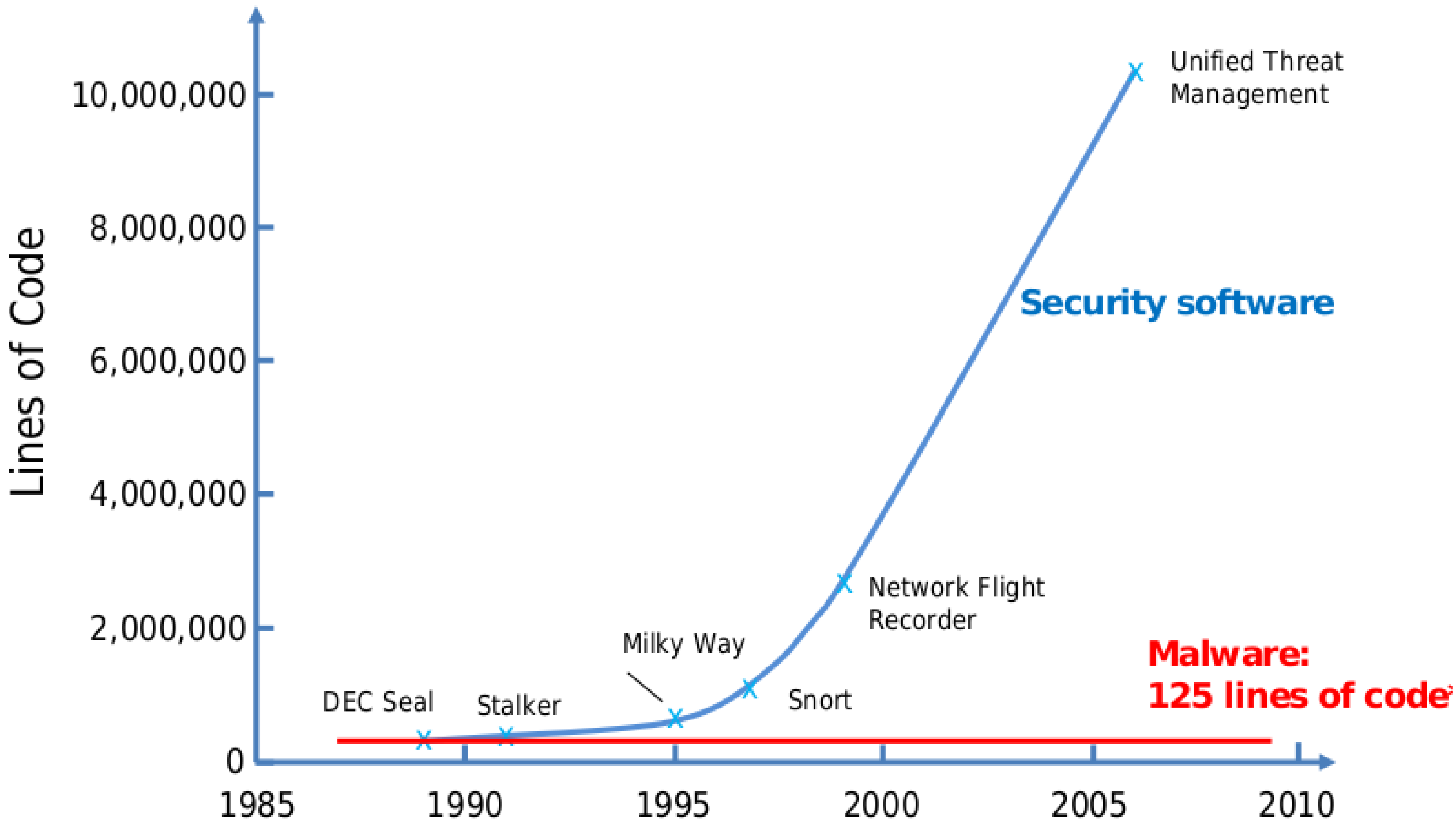
Hardware is now abundant



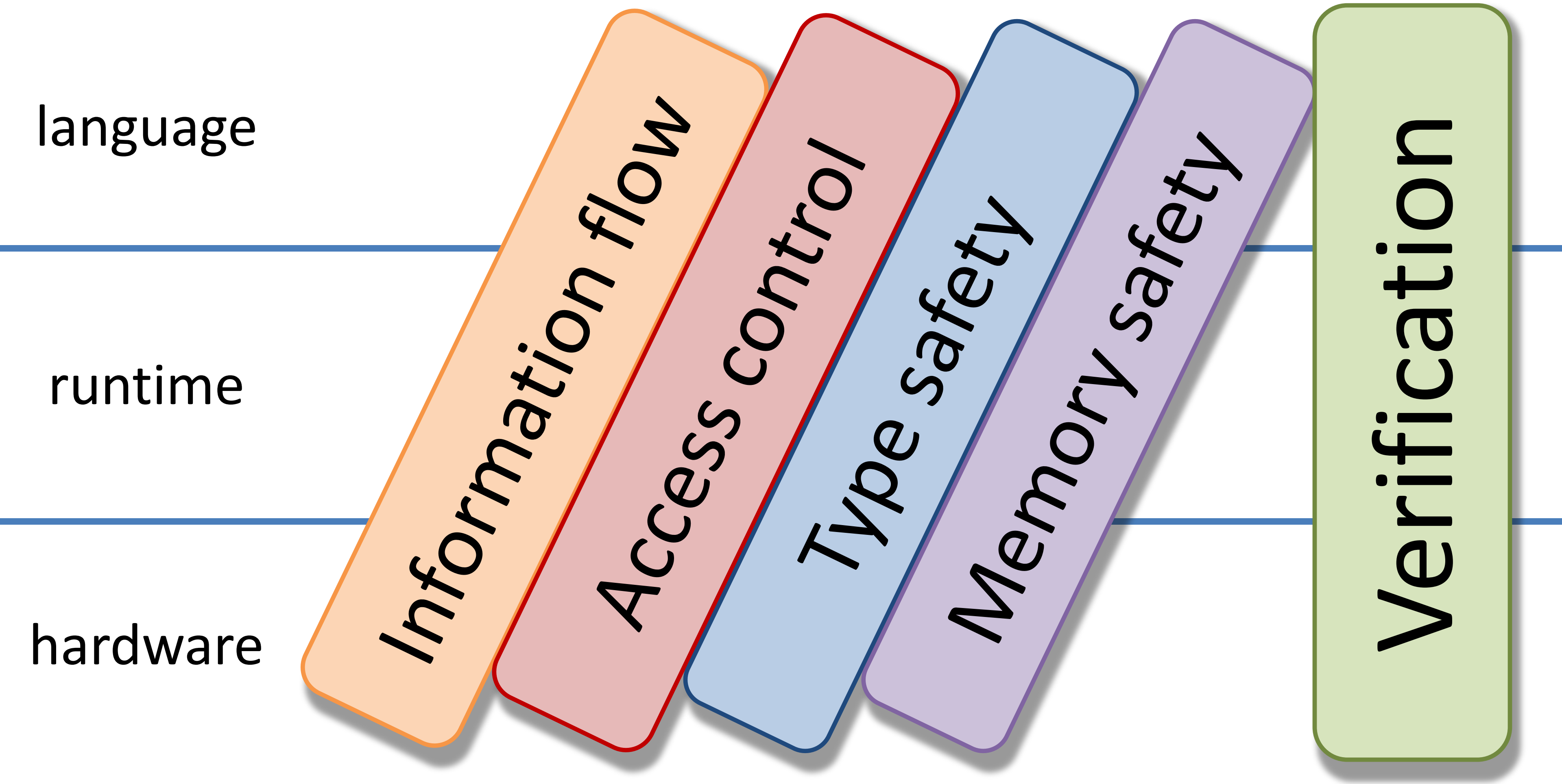
Formal methods are better now

- **random testing**
  - QuickCheck [Claessen & Hughes, ICFP’00]
- **automatic theorem provers & SMT solvers**
- **machine-checked proofs**
  - CompCert [Leroy, POPL’06]
  - seL4 [Klein et al, SOSP’09]
  - CertiCrypt [Barthe et al., POPL’09]
  - ZKCrypt [Almeida et al, CCS’12]

Security is much more important



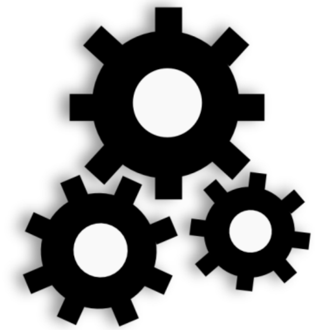
Time for a redesign!



Language (Breeze)

- testing ground for ideas we port to lower levels
- **type and memory safe** high-level language
  - **dynamically typed +** dynamically-checked contracts
- **functional core** (λ) + state(!) + concurrency (π)
  - message-passing communication (channels)
- built-in **fine-grained protection mechanisms**:
  - values are attached **security labels** (e.g. public/secret)
  - **dynamic information flow control** (IFC)
  - **discretionary access control** (clearance)

Runtime system

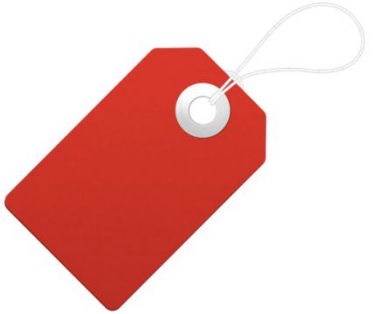


- manages:
  - **time**: scheduler
  - **memory**: allocator, garbage collector
  - **communication and resources**: channels
  - **protection**: principals, authorities, and tags (PAT)
- small trusted computing base
- comparimentalized
  - a dozen mutually distrustful servers (least privilege)

Hardware

- all instructions have well-defined semantics
  - abstractions strictly enforced
- low-**fat pointers**
  - can’t access/write out of frame bounds
- **dynamic types**
  - can’t turn ints into pointers (unforgeable **capabilities**)
- **authority** + **closures/gates** (λ) + **protected stack**
  - fine-grained privilege separation
- programmable **tag management unit** (TMU)

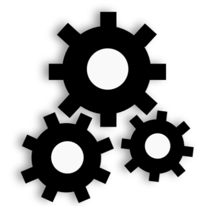
Tag management



- **every word tagged** with arbitrary pointer
  - only runtime system interprets these pointers
- on **each instruction** TMU looks up tags of operands in a **hardware rule cache**
  - found → rule provides tags on results (no delay)
  - not found → trap to software (PAT server)
- **access control + IFC** enforced at lowest level

Project status (2/4 years)

- **language**:
  - stable interpreter, work-in-progress compiler
  - applications: e.g. web server running wiki
  - Coq proofs for various core calculi (non-interference)
- **runtime**:
  - detailed design, some prototype servers
  - work on testing+verifying simplified PAT server



- **hardware**:
  - full-fledged un-optimized FPGA prototype
  - novel instruction set, simulators, debugger, ...
  - executable instruction set semantics in Coq

MY RESEARCH

Pre-SAFE work

- **crypto protocols**
  - tools aiding *design*, *analysis*, and *implementation*
  - **more expressive type systems** (e.g. first one for ZK) [CCS’08, CSF’09, TOSCA’11, PhD thesis]
  - **remote electronic voting** [CSF’08]
  - **code generation** [NFM’12]
- **data processing language (Microsoft “M”)**
  - **semantic subtyping** [ICFP’10, JFP’12]
  - **verification condition generation** [CPP’11]

SAFE work

All Your IFCException Are Belong To Us

Robust Exception Handling for Sound Fine-Grained Dynamic IFC

joint work with Michael Greenberg, Ben Karel,
Benjamin Pierce, and Greg Morrisett

Exception handling

- we wanted all Breeze errors to be **recoverable**
  - **including IFC violations**
- however, existing work assumes errors are **fatal**
  - makes some things easier ... at the expense of others

+secrecy **+integrity** **–availability**

Problem #1: IFC exceptions reveal information about labels

- labels are themselves information channels
- get soundness by preventing secrets from leaking either ***into*** or ***out of*** label channel

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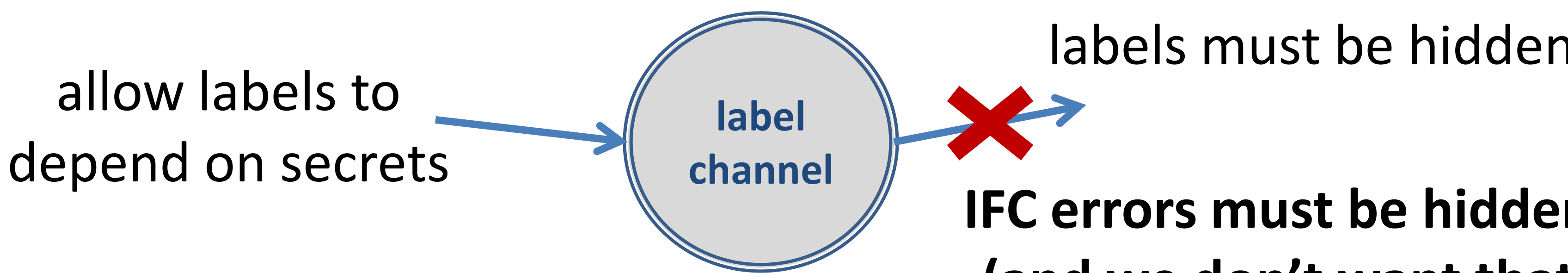


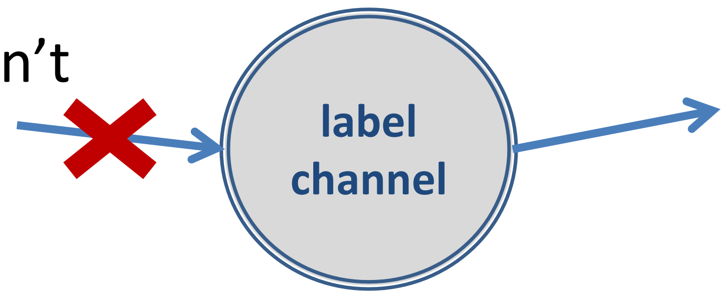
IFC errors must be hidden!
(and we don’t want that)

```
if h@secret then ()@secret else ()@top-secret
```

Problem #1: IFC exceptions reveal information about labels

- labels are themselves information channels
- get soundness by preventing secrets from leaking ~~either~~ ***into*** ~~or **_out of_**~~ label channel

enforce that labels don’t depend on secrets



labels and IFC errors can be observed

Solution #1: brackets

```
top-secret[if h@secret then ()@secret else ()@top-secret]
```

Problem #2: exceptions destroy control flow join points

- ending brackets have to be control flow join points
  - ```
    try
      let _ = secret[if h then throw Ex] in
      false
    catch Ex => true
    ```
- brackets need to delay all exceptions!
  - `secret[if true@secret then throw Ex] => “(Error Ex)@secret”`
  - `secret [if false@secret then throw Ex] => “(Success ())@secret”`
- similarly for failed brackets
  - `secret[42@top-secret] => “(Error EBracket)@secret”`
    ```

Solution #2: Delayed exceptions

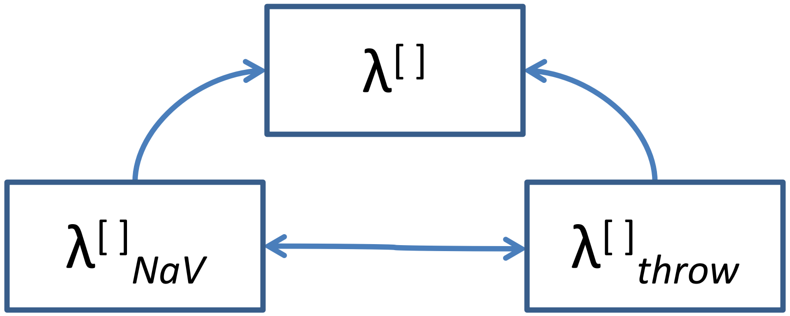
- **delayed exceptions unavoidable**
  - still have a choice how to propagate them
- we studied **two alternatives** for error handling:
  1. **mix active and delayed exceptions** ($\lambda^{[\,]}_{throw}$)
  2. **only delayed exceptions** ($\lambda^{[\,]}_{NaV}$)
     - delayed exception = not-a-value (NaV)
     - NaVs are first-class replacement for values
     - NaVs propagated solely via data flow
     - NaVs are labeled and pervasive
     - more radical solution; implemented by Breeze

What’s in a NaV?

- error message
  - `EDivisionByZero (“can’t divide %1 by 0”, 42)
- stack trace
  - pinpoints error **origin**
    (not the billion-dollar mistake)
- propagation trace
  - how did the error make it here?

Formal results

- proved termination-insensitive **non-interference** in Coq for $\lambda^{[\,]}$, $\lambda^{[\,]}_{NaV}$, and $\lambda^{[\,]}_{throw}$
  - for $\lambda^{[\,]}_{NaV}$ even with all debugging aids; **error-sensitive**
- in our setting NaVs and catchable exceptions have **equivalent expressive power**
  - translations validated by QuickChecking extracted code



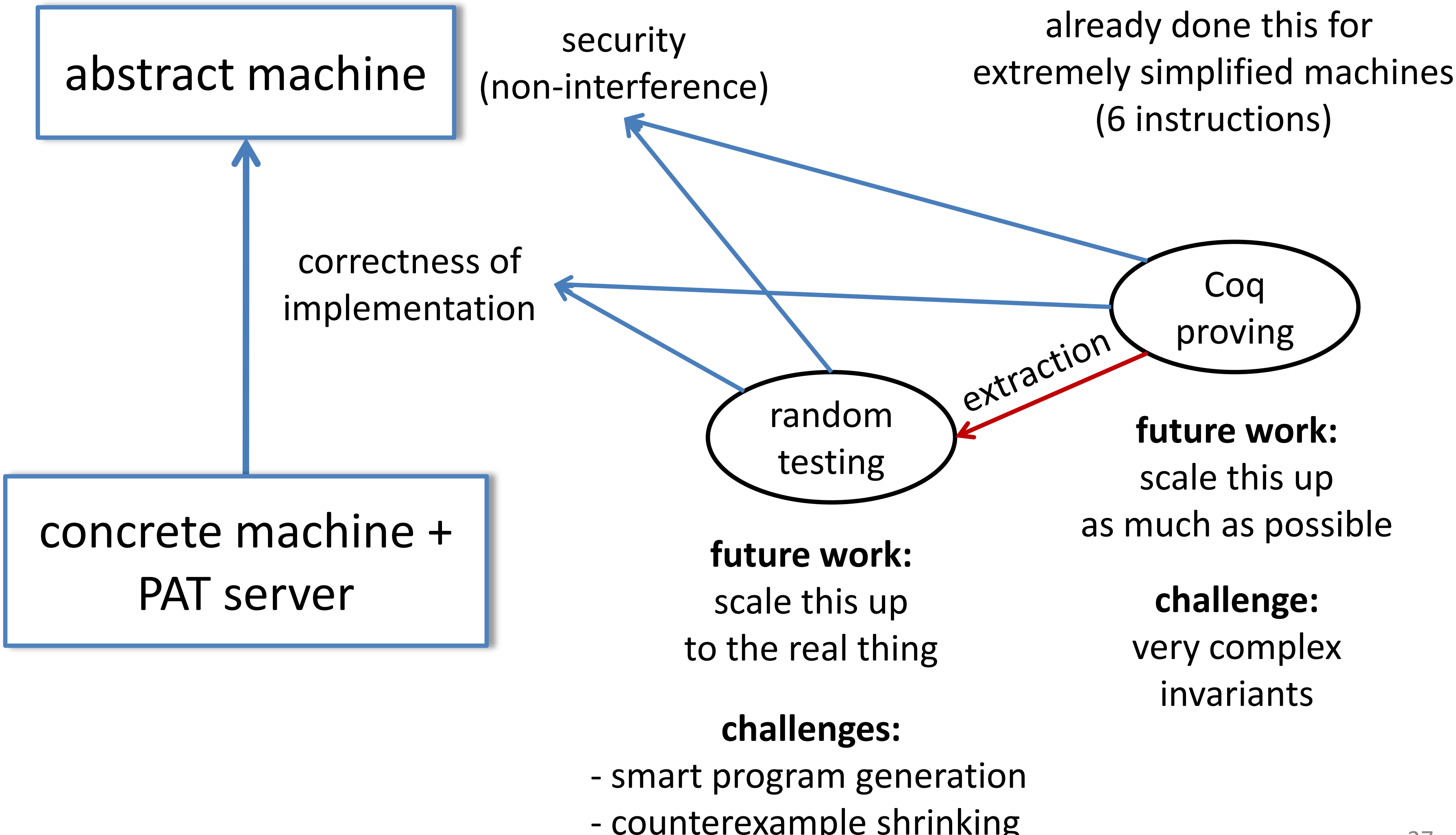
Summary for IFC exceptions

- reliable error handling *possible* even for sound fine-grained dynamic IFC systems
- we study two mechanisms ($\lambda^{[\,]}_{NaV}$ and $\lambda^{[\,]}_{throw}$)
  - **all errors recoverable**, even IFC violations
  - key ingredients:
    sound public labels (brackets) + **delayed exceptions**
  - quite radical design (not backwards compatible!)
- gathering practical experience with NaVs:
  - issues are surmountable
  - writing good error recovery code is still hard

Ongoing work

- **testing and verifying the PAT server**
- protecting data integrity with signature labels
- implementing Breeze labels cryptography

Testing and verifying PAT server



Post-SAFE work?

- **software-hardware co-design for security-critical high-assurance devices**
  – electronic voting, driver assistance, medical devices
    - limited/fixed functionality
    - security more important than backwards compatibility
  – existing devices often blatantly vulnerable
  – making security analysis part of design process
  – focus more on research (compared to CRASH/SAFE)
- fine-grained access control and integrity protection for mobile devices

Possible collaborations at TU Darmstadt

- Prof. **Heiko Mantel** (dynamic IFC and concurrency)
- Prof. **Ahmad-Reza Sadeghi** (smartphone or automobile security),
- Prof. **Melanie Volkamer** (remote electronic voting),
- Dr. **Thomas Schneider** (formal proofs for SMPC & ZK compilers),
- Dr. **Eric Bodden** (security monitoring for mobile devices)
- Prof. **Thomas Streicher** (logics and semantics)

THE END

BACKUP SLIDES

Sound dynamic IFC possible

- Non-interference can be obtained purely dynamically!
  - [Krohn & Tromer, 2009], [Sabelfeld & Russo, 2009], [Austin & Flanagan, 2009]
- Preventing implicit flows:

```
let lref = ref low false in
if h then
  lref := true;
lref := false
...
```

pc=high

potential bad flow -> halt program

false alarm (program non-interferent)

- Even functional code can leak via control flow:
  - `if h then true else false`
  - semantics of conditional:
    - **if** true@high **then** true **else** false => true@high